

## REGULARITY OF CUBIC GRAPH WITH APPLICATION

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**Abstract.** A cubic graph is a generalized structure of a fuzzy graph that gives more precision, flexibility and compatibility to a system when compared with systems that are designed using fuzzy graphs. In this paper, some properties of an edge regular cubic graph are given. Particularly, strongly regular, edge regular and bi-regular cubic graphs are defined and the necessary and sufficient condition for a cubic graph to be strongly regular is studied. Likewise, we have introduced a partially edge regular cubic graph and fully edge regular cubic graph with suitable illustrations. Finally, we gave an application of cubic digraph in travel time.

*Key words and Phrases:* Cubic graph, strongly regular cubic graph, bi-regular cubic graph

**Abstrak.** Graf kubik merupakan suatu struktur perumuman dari graf fuzzy yang memberikan banyak presisi, fleksibilitas, dan kompatibilitas terhadap suatu sistem yang dirancang menggunakan graf fuzzy. Dalam paper ini, dipelajari beberapa sifat dari suatu graf kubik yang teratur sisi. Didefinisikan pula graf kubik teratur kuat yang teratur sisi dan *bi-regular* serta dipelajari syarat perlu dan syarat cukup bagi suatu graf kubik agar menjadi graf yang teratur kuat. Kemudian, diperkenalkan graf kubik yang teratur sisi secara parsial dan graf kubik teratur sisi secara penuh dengan beberapa ilustrasi diberikan. Terakhir, diberikan aplikasi dari graf kubik pada *travel time*.

*Kata kunci:* Graf kubik, Graf kubik teratur kuat, Graf kubik *bi-regular*.

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2000 Mathematics Subject Classification: 05C72, 05C99.

Received: 25-08-2017, revised : 20-05-2018, accepted: 31-05-2018.

## 1. INTRODUCTION

Zadeh introduced the concept of fuzzy set in his seminal paper [22] of 1965. A fuzzy set of a universe  $X$  is a function from  $X$  into the unit closed interval  $[0, 1]$  of real number. In [23] Zadeh made an extension of the concept of a fuzzy set by an interval-valued fuzzy set, i.e., a fuzzy set with an interval-valued membership function. The first definition of fuzzy graphs was proposed by Kaufmann [5] in 1973. Interval-valued fuzzy sets have been actively used in real-life applications. For example, Sambuc [14] in medical diagnosis in thyroidian pathology, Kohout [7] also in medicine, Turksen in preferences modeling [21], etc. These works and others showed the importance of these sets. Jun et al. [4] introduced cubic sets. The fuzzy graph theory as a generalization of Euler's graph theory was first introduced by Rosenfeld [13] in 1975. Later, Bhattacharya [2] gave some remarks on fuzzy graphs and some operations on fuzzy graphs were introduced by Mordeson and Peng [8]. The complement of a fuzzy graph was defined by Mordeson [9] and further studied by Sunitha and Vijayakumar [15]. Hongmei and Lianhua gave the definition of interval-valued fuzzy graphs [3]. Akram and Dudek defined some operations on interval-valued fuzzy graphs [1]. Rashmanlou et al. [10, 11, 12] introduced some properties of highly irregular interval-valued fuzzy graphs, and new concepts of bipolar fuzzy graphs. Karunambigai et al. [6] introduced edge regular intuitionistic fuzzy graph. Samanta and Pal [15, 16, 17, 18, 19, 20] defined fuzzy tolerance graph, fuzzy threshold graph, fuzzy  $k$ -competition graph and  $p$ -competition fuzzy graph and new concepts of fuzzy planar graph. The major role of cubic graph theory in computer applications is the development of graph algorithms. These algorithms are used to solve problems that are modeled in the form of graphs and the corresponding computer science application problems. One of the most widely studied classes of cubic graphs is regular cubic graphs. Theoretical concepts of cubic graphs are highly utilized by computer science applications. Especially in research areas of computer science such as data mining, image segmentation, clustering, image capturing and networking. The cubic graphs are more flexible and compatible than fuzzy graphs due to the fact that they have many applications in networks. They show up in many contexts. For example,  $r$ -regular cubic graphs with connectivity and edge-connectivity equal to  $r$  play a key role in designing reliable communication networks. Hence, in this paper some properties of an edge regular cubic graph are given. Particularly, strongly regular, edge regular and biregular cubic graphs are defined and the necessary and sufficient condition for a cubic graph to be strongly regular is studied. Also, we have introduced a partially edge regular cubic graph and fully edge regular cubic graph with suitable illustrations. After introductory Section 1, some basic definitions are given in Section 2. In Section 3, the concepts of regularity of cubic graphs are defined. In Section 4, an application is given in travel time. At last conclusion is given in Section 5.

## 2. PRELIMINARIES

A graph is an ordered pair  $G = (V, E)$ , where  $V$  is the set of vertices of  $G$  and  $E$  is the set of edges of  $G$ . A subgraph of a graph  $G = (V, E)$  is a graph  $H = (W, F)$ , where  $W \subseteq V$  and  $F \subseteq E$ . A fuzzy graph  $G = (\sigma, \mu)$  is a pair of functions  $\sigma : V \rightarrow [0, 1]$  and  $\mu : V \times V \rightarrow [0, 1]$  with  $\mu(u, v) \leq \sigma(u) \wedge \sigma(v)$ , for all  $u, v \in V$ , where  $V$  is a finite non-empty set and  $\wedge$  denote minimum. We introduce below necessary notions and present

a few auxiliary results that will be used throughout the paper.

A map  $\lambda : X \rightarrow [0, 1]$  is called a fuzzy subset of  $X$ . For any two fuzzy subsets  $\lambda$  and  $\mu$  of  $X$ ,  $\lambda \subseteq \mu$  means that, for all  $x \in X$ ,  $\lambda(x) \leq \mu(x)$ . The symbol  $\lambda \wedge \mu$  and  $\lambda \vee \mu$  will mean the following fuzzy subsets of  $X$ .

$$(\lambda \wedge \mu)(x) = \lambda(x) \wedge \mu(x) \text{ and } (\lambda \vee \mu)(x) = \lambda(x) \vee \mu(x), \text{ for all } x \in X.$$

Let  $X$  be a non-empty set. A function  $A : X \rightarrow [I]$  is called an interval-valued fuzzy set (shortly, an IVF set) in  $X$ . Let  $[I]^X$  stands for the set of all IVF sets in  $X$ . For every  $A \in [I]^X$  and  $x \in X$ ,  $A(x) = [A^-(x), A^+(x)]$  is called the degree of membership of an element  $x$  to  $A$ , where  $A^- : X \rightarrow I$  and  $A^+ : X \rightarrow I$  are fuzzy sets in  $X$  which are called a lower fuzzy set and an upper fuzzy set in  $X$ , respectively. For simplicity, we denote  $A = [A^-, A^+]$ . For every  $A, B \in [I]^X$ , we define  $A \subseteq B$  if and only if  $A(x) \leq B(x)$ , for all  $x \in X$ .

**Definition 2.1.** Let  $A = [A^-, A^+]$ , and  $B = [B^-, B^+]$  be two interval-valued fuzzy set in  $X$ . Then we define  $\text{rmin}\{A(x), B(x)\} = [\min\{A^-(x), B^-(x)\}, \min\{A^+(x), B^+(x)\}]$ ,  $\text{rmax}\{A(x), B(x)\} = [\max\{A^-(x), B^-(x)\}, \max\{A^+(x), B^+(x)\}]$ .

**Definition 2.2.** Let  $X$  be a non-empty set. By a cubic set in  $X$ , we mean a structure  $A = \langle x, A(x), \lambda(x) : x \in X \rangle$  in which  $A$  is an interval-valued fuzzy sets in  $X$  and  $\lambda$  is a fuzzy set in  $X$ . A cubic set  $A = \langle x, A(x), \lambda(x) : x \in X \rangle$  is simply denoted by  $A = \langle A, \lambda \rangle$ . The collection of all cubic sets in  $X$  is denoted by  $CP(X)$ .

**Definition 2.3.** A cubic graph is a triple  $G = (G^*, P, Q)$  where  $G^* = (V, E)$  is a graph,  $P = (\widetilde{\mu}_P, \lambda_P)$  is a cubic set on  $V$  and  $Q = (\widetilde{\mu}_Q, \widetilde{\lambda}_Q)$  is a cubic set on  $V \times V$  such that  $\widetilde{\mu}_Q(xy) \leq \text{rmin}\{\widetilde{\mu}_P(x), \widetilde{\mu}_P(y)\}$  and  $\widetilde{\lambda}_Q(xy) \geq \text{rmax}\{\lambda_P(x), \lambda_P(y)\}$

The underlying crisp graph of a cubic graph  $G = (A, B)$ , is the graph  $G = (V, E)$ , where  $V = \{v : \widetilde{\mu}_P(v) > 0 \text{ and } \lambda_P(v) > 0\}$  and  $E = \{\{u, v\} : \widetilde{\mu}_Q(\{u, v\}) > 0, \widetilde{\lambda}_Q(\{u, v\}) > 0\}$ .  $V$  is called the vertex set and  $E$  is called the edge set. A cubic graph maybe also denoted as  $G = (V, E)$ .

**Definition 2.4.** A cubic graph  $G = (G^*, P, Q)$  is called complete if  $\widetilde{\mu}_Q(xy) = \text{rmin}\{\mu_P(x), \mu_P(y)\}$  and  $\widetilde{\lambda}_Q(xy) = \text{rmax}\{\lambda_P(x), \lambda_P(y)\}$  for all  $x, y \in V$ .

**Definition 2.5.** A cubic graph  $G = (G^*, P, Q)$  is called strong if  $\widetilde{\mu}_Q(xy) = \text{rmin}\{\mu_P(x), \mu_P(y)\}$  and  $\widetilde{\lambda}_Q(xy) = \text{rmax}\{\lambda_P(x), \lambda_P(y)\}$  for all  $xy \in E$ .

**Definition 2.6.** The complement of a cubic graph  $G = (A, B)$  is a cubic graph  $\overline{G} = (\overline{G}^*, \overline{P}, \overline{Q})$ , where  $\overline{P} = (\overline{\mu}_P, \overline{\lambda}_P)$  and  $\overline{Q} = (\overline{\mu}_Q, \overline{\lambda}_Q)$  is defined by:  $\overline{\mu}_Q(xy) = \text{rmin}\{\mu_P(x), \mu_P(y)\} - \widetilde{\mu}_Q(xy)$  and  $\overline{\lambda}_Q(xy) = \widetilde{\lambda}_Q(xy) - \text{rmax}\{\lambda_P(x), \lambda_P(y)\}$

**Definition 2.7.** Let  $G = (V, E)$  be a cubic graph.

(i) The neighborhood degree of a vertex  $v$  is defined as  $D_N(v) = (d_{N_{\overline{\mu}_P}}(v), d_{N_{\lambda_Q}}(v))$ , where  $d_{N_{\mu_P}}(v) = \left( \sum_{w \in N(\overline{\mu}_P)} \overline{\mu}_P(w), \sum_{w \in N(\mu_P^+)} \mu_P^+(w) \right)$  and  $d_{N_{\lambda_Q}}(v) = \sum_{w \in N(\lambda_Q)} \lambda_Q(w)$ .

(ii) The degree of a vertex  $v_i$  is defined by  $d_G(v_i) = (d_{\overline{\mu}_P}(v_i), d_{\lambda_Q}(v_i)) = (k_1, k_2)$ , where

$$k_1 = d_{\widetilde{\mu}_P}(v_i) = \left( \sum_{v_i \neq v_j} \widetilde{\mu}_Q(v_i v_j), \sum_{v_i \neq v_j} \widetilde{\mu}_Q^+(v_i v_j) \right)$$

$$\text{and } k_2 = d_{\lambda_Q}(v_i) = \sum_{v_i \neq v_j} \lambda_Q(v_i v_j).$$

**Definition 2.8.** A cubic graph  $G = (V, E)$  is said to be

(i)  $(k_1, k_2)$ -regular if  $d_G(v_i) = (k_1, k_2)$ , for all  $v_i \in V$  and also  $G$  is said to be a regular cubic graph of degree  $(k_1, k_2)$ .

(ii) bipartite if the vertex set  $V$  can be partitioned into two non-empty sets  $V_1$  and  $V_2$  such that

(a)  $\widetilde{\mu}_Q(v_i v_j) = 0$  and  $\lambda_Q(v_i v_j) = 0$ , if  $(v_i, v_j) \in V_1$  or  $(v_i, v_j) \in V_2$

(b)  $\widetilde{\mu}_Q(v_i v_j) = 0$ ,  $\lambda_Q(v_i v_j) > 0$ , if  $v_i \in V_1$  or  $v_j \in V_2$

(c)  $\widetilde{\mu}_Q(v_i v_j) > 0$ ,  $\lambda_Q(v_i v_j) = 0$ , if  $v_i \in V_1$  or  $v_j \in V_2$ , for some  $i$  and  $j$ .

**Definition 2.9.** Let  $G^* = (V, E)$  be a crisp graph and let  $e = v_i v_j$  be an edge in  $G^*$ . Then, the degree of an edge  $e = v_i v_j \in E$  is defined as  $d_{G^*}(v_i v_j) = d_{G^*}(v_i) + d_{G^*}(v_j) - 2$ .

**Definition 2.10.** (i) The order of  $G$  is defined to be  $O(G) = (O_{\widetilde{\mu}_P}, O_{\lambda_P})$ , where  $O_{\widetilde{\mu}_P} = \sum_{u \in V} \widetilde{\mu}_P(u)$  and  $O_{\lambda_P} = \sum_{u \in V} \lambda_P(u)$ .

(ii) The size of  $G$  is defined to be

$$S(G) = (S_{\widetilde{\mu}_Q}(G), S_{\lambda_Q}(G)), \text{ where } S_{\widetilde{\mu}_Q}(G) = \sum_{u \neq v} \widetilde{\mu}_Q(uv) \text{ and } S_{\lambda_Q}(G) = \sum_{u \neq v} \lambda_Q(uv).$$

### 3. ISOMORPHIC PROPERTIES OF NEIGHBORLY IRREGULAR AND HIGHLY IRREGULAR CUBIC GRAPHS

In this section, we define weak isomorphism, co-weak isomorphism and isomorphism of neighborly irregular cubic graphs and prove that.

**Definition 3.1.** Let  $G = (V, E)$  be a cubic graph.

(i) The degree of an edge  $e_{ij} \in E$  is defined as

$$d_{\widetilde{\mu}_Q}(e_{ij}) = d_{\widetilde{\mu}_P}(v_i) + d_{\widetilde{\mu}_P}(v_j) - 2\widetilde{\mu}_Q(v_i v_j) \text{ or } d_{\widetilde{\mu}_Q}(e_{ij}) = \sum_{\substack{v_i v_k \in E \\ k \neq j}} \widetilde{\mu}_Q(v_i v_k) + \sum_{\substack{v_k v_j \in E \\ k \neq i}} \widetilde{\mu}_Q(v_k v_j).$$

$$d_{\lambda_Q}(e_{ij}) = d_{\lambda_P}(v_i) + d_{\lambda_P}(v_j) - 2\lambda_Q(v_i v_j) \text{ or } d_{\lambda_Q}(e_{ij}) = \sum_{\substack{v_i v_k \in E \\ k \neq j}} \lambda_Q(v_i v_k) + \sum_{\substack{v_k v_j \in E \\ k \neq i}} \lambda_Q(v_k v_j).$$

(ii) The minimum edge degree of  $G$  is  $\delta_E(G) = (\delta_{\widetilde{\mu}_Q}(G), \delta_{\lambda_Q}(G))$ , where  $\delta_{\widetilde{\mu}_Q}(G) = \wedge \{d_{\widetilde{\mu}_Q}(e_{ij}) \mid e_{ij} \in E\}$  and  $\delta_{\lambda_Q}(G) = \wedge \{d_{\lambda_Q}(e_{ij}) \mid e_{ij} \in E\}$ .

(iii) The maximum edge degree of  $G$  is  $\Delta_E(G) = (\Delta_{\widetilde{\mu}_Q}(G), \Delta_{\lambda_Q}(G))$ , where  $\Delta_{\widetilde{\mu}_P}(G) = \vee \{d_{\widetilde{\mu}_Q}(e_{ij}) \mid e_{ij} \in E\}$  and  $\Delta_{\lambda_Q}(G) = \vee \{d_{\lambda_Q}(e_{ij}) \mid e_{ij} \in E\}$ .

(iv) The total edge degree of an edge  $e_{ij} \in E$  is defined as

$$td_{\widetilde{\mu}_Q}(e_{ij}) = \sum_{\substack{v_i v_k \in E \\ k \neq j}} \widetilde{\mu}_Q(v_i v_k) + \sum_{\substack{v_k v_j \in E \\ k \neq i}} \widetilde{\mu}_Q(v_k v_j) + \widetilde{\mu}_Q(e_{ij}), \quad td_{\lambda_Q}(e_{ij}) = \sum_{\substack{v_i v_k \in E \\ k \neq j}} \lambda_Q(v_i v_k) + \sum_{\substack{v_k v_j \in E \\ k \neq i}} \lambda_Q(v_k v_j) +$$

$$\lambda_Q(e_{ij}).$$

(v) The edge degree of  $G$  is defined by  $d_G(e_{ij}) = (d_{\widetilde{\mu}_Q}(e_{ij}), d_{\lambda_Q}(e_{ij}))$  and the total edge degree of  $G$  is defined by  $td_G(e_{ij}) = (td_{\widetilde{\mu}_Q}(e_{ij}), td_{\lambda_Q}(e_{ij}))$ .

**Example 3.1.** Consider the cubic graph  $G = (V, E)$  in Figure 3.1, where  $V = \{a, b, c, d\}$  and  $E = \{ab, ad, bc, bd, cd\}$ . Then

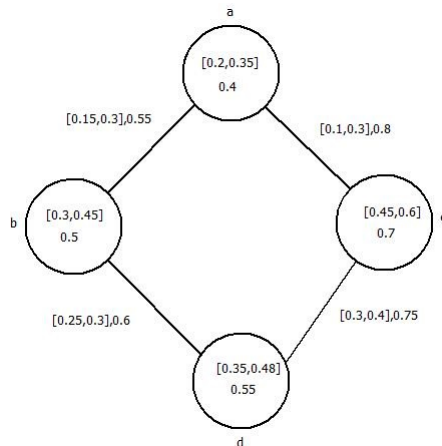


FIGURE 1. Cubic graph

$d_{\mu_Q^-}(ac) = [0.7, 1.05]$ ,  $d_{\lambda_Q}(ac) = 1.1$ ,  $d_G(uw) = [0.7, 1.05]$ ,  $1.1$ ,  $td_{\mu_Q^-}(ac) = [0.7 + 0.1 = 0.8, 1.05 + 0.25 = 1.30]$  and  $td_{\lambda_Q}(ac) = 1.7 + 0.4 = 2.1$ . Hence,  $td_G(ac) = [0.8, 1.30], 2.1$  ( $e_{ij} = (u_i, u_j)$ ).

**Definition 3.2.** Let  $G = (V, E)$  be a cubic graph.

(i) If each edge in  $G$  has the same degree  $(l_1, l_2)$ , then  $G$  is said to be an edge regular cubic graph.

(ii) If each edge in  $G$  has the same total degree  $(t_1, t_2)$ , then  $G$  is said to be a totally edge regular cubic graph.

**Example 3.2.** Consider the cubic graph  $G = (V, E)$  as in Figure 2, where  $V = \{u_1, u_2, u_3, u_4\}$  and  $E = \{u_1u_2, u_1u_3, u_3u_4, u_2u_4\}$ . Then

$$d_G(e_{12}) = d_G(e_{24}) = d_G(e_{34}) = d_G(e_{13}) = [0.1, 0.3], 1.1.$$

**Theorem 3.1.** Let  $G = (V, E)$  be a cubic graph on a cycle  $G^*$ . Then

$$\sum_{v_i \in V} d_G(v_i) = \sum_{v_i v_j \in E} d_G(v_i v_j)$$

*Proof.* Let  $G = (V, E)$  be a cubic graph and  $G^*$  be a cycle  $v_1 v_2 v_3 \cdots v_n v_1$ . Then

$$\sum_{i=1}^n d_G(v_i v_{i+1}) = \left( \sum_{i=1}^n d_{\mu_Q^-}(v_i v_{i+1}), \sum_{i=1}^n d_{\lambda_Q}(v_i v_{i+1}) \right). \text{ Also } \widetilde{\mu}_P = \{\mu_P^-, \mu_P^+\} \text{ and } \widetilde{\mu}_Q = \{\mu_Q^-, \mu_Q^+\}$$

Now we have

$$\begin{aligned} & \sum_{i=1}^n d_{\mu_Q^-}(v_i v_{i+1}) = \\ & d_{\mu_Q^-}(v_1 v_2) + d_{\mu_Q^-}(v_2 v_3) + \cdots + d_{\mu_Q^-}(v_n v_1), \text{ where } v_{n+1} = v_1 \\ & = d_{\mu_P^-}(v_1) + d_{\mu_P^-}(v_2) - 2\widetilde{\mu}_Q(v_1 v_2) + d_{\mu_P^-}(v_2) + d_{\mu_P^-}(v_3) \\ & = 2\widetilde{\mu}_P(v_2 v_3) + \cdots + d_{\mu_P^-}(v_n) + d_{\mu_P^-}(v_1) - 2\widetilde{\mu}_Q(v_n v_1) \\ & = 2d_{\mu_P^-}(v_1) + 2d_{\mu_P^-}(v_2) + \cdots + 2d_{\mu_P^-}(v_n) \end{aligned}$$

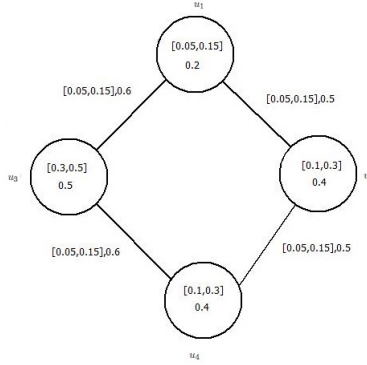


FIGURE 2.  $[0.1,0.3],1.1$  - edge regular cubic graph

$$\begin{aligned}
 &= 2(\widetilde{\mu}_Q(v_1v_2) + \widetilde{\mu}_Q(v_2v_3) + \cdots + \widetilde{\mu}_Q(v_nv_1)) \\
 &= 2 \sum_{v_i \in V} d_{\widetilde{\mu}_P}(v_i) \\
 &\quad - 2 \sum_{i=1}^n \widetilde{\mu}_Q(v_i v_{i+1}) \\
 &= \sum_{v_i \in V} d_{\widetilde{\mu}_P}(v_i) + 2 \sum_{i=1}^n \widetilde{\mu}_Q(v_i v_{i+1}) \\
 &\quad - 2 \sum_{i=1}^n \widetilde{\mu}_Q(v_i v_{i+1}) \\
 &= \sum_{v_i \in V} \widetilde{\mu}_P(v_i).
 \end{aligned}$$

$$\text{Similarly, } \sum_{i=1}^n d_{\lambda_Q}(v_i v_{i+1}) = \sum_{v_i \in V} d_{\lambda_P}(v_i).$$

Hence,  $\sum_{i=1}^n d_G(v_i v_{i+1}) = \left( \sum_{v_i \in V} d_{\widetilde{\mu}_P}(v_i), \right.$

$$\begin{aligned}
 &\left. \sum_{v_i \in V} d_{\lambda_P}(v_i) \right) \\
 &= \sum_{v_i \in V} d_G(v_i). \quad \square
 \end{aligned}$$

**Remark 3.1.** Let  $G = (V, E)$  be a cubic graph on a crisp graph  $G^*$ . Then,  $\sum_{v_i, v_j \in E} d_G(v_i v_j) =$

$$\left( \sum_{v_i, v_j \in E} d_{G^*}(v_i v_j) \widetilde{\mu}_Q(v_i v_j), \right.$$

$\left. \sum_{v_i, v_j \in E} d_{G^*}(v_i v_j) \lambda_Q(v_i v_j) \right)$ , where  $d_{G^*}(v_i v_j) = d_{G^*}(v_i) + d_{G^*}(v_j) - 2$ , for all  $v_i v_j \in E$ .

**Theorem 3.2.** Let  $G = (V, E)$  be a cubic graph on a  $k$ -regular crisp graph  $G^*$ . Then,

$$\sum_{v_i, v_j \in E} d_G(v_i v_j) = \left( (k-1) \sum_{v_i \in V} d_{\widetilde{\mu}_P}(v_i), (k-1) \sum_{v_i \in V} d_{\lambda_P}(v_i) \right).$$

*Proof.* By Remark 3.6, we have  $\sum_{v_i, v_j \in E} d_G(v_i v_j)$

$$\begin{aligned}
 &= \left( \sum_{v_i, v_j \in E} d_{G^*}(v_i v_j) \widetilde{\mu}_Q(v_i v_j), \sum_{v_i, v_j \in E} \right. \\
 &\left. d_{G^*}(v_i v_j) \lambda_Q(v_i v_j) \right)
 \end{aligned}$$

$$= \left( \sum_{v_i, v_j \in E} (d_{G^*}(v_i) + d_{G^*}(v_j) - 2) \widetilde{\mu}_Q(v_i v_j), \right.$$

$\left. \sum_{v_i, v_j \in E} (d_{G^*}(v_i) + d_{G^*}(v_j) - 2) \lambda_Q(v_i v_j) \right)$ . Since  $G^*$  is a regular crisp graph,  $d_{G^*}(v_i) = k$ ,

for all  $v_i \in V$  and so we have  $\sum_{v_i, v_j \in E} d_G(v_i v_j) = \left( (k+k-2) \sum_{v_i, v_j \in E} \widetilde{\mu}_Q(v_i v_j), (k+k-2) \sum_{v_i, v_j \in E} \lambda_Q(v_i v_j) \right)$ ,

$$\sum_{v_i, v_j \in E} d_G(v_i v_j) = \left( 2(k-1) \sum_{v_i, v_j \in E} \widetilde{\mu}_Q(v_i v_j), 2(k-1) \sum_{v_i, v_j \in E} \lambda_Q(v_i v_j) \right),$$

$$\sum_{v_i, v_j \in E} d_G(v_i v_j) = \left( (k-1) \sum_{v_i \in V} d_{\widetilde{\mu}_P}(v_i), (k-1) \sum_{v_i \in V} d_{\lambda_P}(v_i) \right). \quad \square$$

**Theorem 3.3.** Let  $G = (V, E)$  be a cubic graph on a crisp graph  $G^*$ . Then,  $\sum_{v_i v_j \in E} td_G(v_i v_j) = \left( \sum_{v_i v_j \in E} d_{G^*}(v_i v_j) \widetilde{\mu}_Q(v_i v_j) + \sum_{v_i v_j \in E} \widetilde{\mu}_Q(v_i v_j), \sum_{v_i v_j \in E} d_{G^*}(v_i v_j) \lambda_Q(v_i v_j) + \sum_{v_i v_j \in E} \lambda_Q(v_i v_j) \right)$ .

*Proof.* By definition of total edge degree of  $G$ , we have  $\sum_{v_i v_j \in E} td_G(v_i v_j) = \left( \sum_{v_i v_j \in E} td_{\widetilde{\mu}_Q}(v_i v_j), \sum_{v_i v_j \in E} td_{\lambda_Q}(v_i v_j) \right) = \left( \sum_{v_i v_j \in E} (d_{\widetilde{\mu}_Q}(v_i v_j) + \widetilde{\mu}_Q(v_i v_j)), \sum_{v_i v_j \in E} (d_{\lambda_Q}(v_i v_j) + \lambda_Q(v_i v_j)) \right) = \left( \sum_{v_i v_j \in E} d_{\widetilde{\mu}_Q}(v_i v_j) + \sum_{v_i v_j \in E} \widetilde{\mu}_Q(v_i v_j), \sum_{v_i v_j \in E} d_{\lambda_Q}(v_i v_j) + \sum_{v_i v_j \in E} \lambda_Q(v_i v_j) \right)$ .

By Remark 3.6, we get

$$\begin{aligned} & \sum_{v_i v_j \in E} td_G(v_i v_j) \\ &= \left( \sum_{v_i v_j \in E} d_{G^*}(v_i v_j) \widetilde{\mu}_Q(v_i v_j) + \sum_{v_i v_j \in E} \widetilde{\mu}_Q(v_i v_j), \sum_{v_i v_j \in E} d_{G^*}(v_i v_j) \lambda_Q(v_i v_j) + \sum_{v_i v_j \in E} \lambda_Q(v_i v_j) \right). \end{aligned}$$

□

**Theorem 3.4.** Let  $G = (V, E)$  be a cubic graph. Then  $(t_B, f_B)$  is a constant function if and only if the following are equivalent.

- (i)  $G$  is a edge regular cubic graph.
- (ii)  $G$  is totally edge regular cubic graph.

*Proof.* Assume that  $(\widetilde{\mu}_P, \lambda_Q)$  is a constant function. Then  $\widetilde{\mu}_Q(v_i v_j) = c_1$  and  $\lambda_Q(v_i v_j) = c_2$ , for every  $v_i v_j \in E$ , where  $c_1$  and  $c_2$  are constants. Let  $G$  be an  $(l_1, l_2)$ -edge regular cubic graph. Then, for all  $v_i v_j \in E$ ,  $d_G(v_i v_j) = (l_1, l_2)$  and  $td_G(v_i v_j) = (d_{\widetilde{\mu}_P}(v_i v_j) + \widetilde{\mu}_Q(v_i v_j), d_{\lambda_Q}(v_i v_j) + \lambda_Q(v_i v_j)) = (l_1 + c_1, l_2 + c_2)$ , for all  $v_i v_j \in E$ . Then  $G$  is a totally edge regular. Now, let  $G$  be a  $(t_1, t_2)$ -totally edge regular cubic graph. Then  $td_G(v_i v_j) = (t_1, t_2)$ , for all  $v_i v_j \in E$ . So, we have  $td_G(v_i v_j) = (d_{\widetilde{\mu}_Q}(v_i v_j) + \widetilde{\mu}_P(v_i v_j), d_{\lambda_Q}(v_i v_j) + \lambda_Q(v_i v_j)) = (t_1, t_2)$ . Hence,  $(d_{\widetilde{\mu}_Q}(v_i v_j), d_{\lambda_Q}(v_i v_j)) = (t_1 - \widetilde{\mu}_Q(v_i v_j), t_2 - \lambda_Q(v_i v_j)) = (t_1 - c_1, t_2 - c_2)$ . Then,  $G$  is a  $(t_1 - c_1, t_2 - c_2)$  edge regular cubic graph.

Conversely, assume that (i) and (ii) are equivalent. We have to prove that  $(\widetilde{\mu}_P, \lambda_Q)$  is a constant function. Suppose that  $(\widetilde{\mu}_P, \lambda_Q)$  is not a constant function. Then  $\widetilde{\mu}_Q(v_i v_j) \neq \widetilde{\mu}_Q(v_r v_s)$  and  $\lambda_Q(v_i v_j) \neq \lambda_Q(v_r v_s)$  for at least one pair of  $v_i v_j, v_r v_s \in E$ . Let  $G$  be an  $(l_1, l_2)$  edge regular cubic graph. Then,  $d_G(v_i v_j) = d_G(v_r v_s) = (l_1, l_2)$ . Hence for all  $v_i v_j \in E$  and for all  $v_r v_s \in E$ ;  $td_G(v_i v_j) = (d_{\widetilde{\mu}_Q}(v_i v_j) + \widetilde{\mu}_Q(v_i v_j), d_{\lambda_Q}(v_i v_j) + \lambda_Q(v_i v_j)) = (l_1 + \widetilde{\mu}_Q(v_i v_j), l_2 + \lambda_Q(v_i v_j))$   $td_G(v_r v_s) = (\widetilde{\mu}_Q(v_r v_s) + \widetilde{\mu}_Q(v_r v_s), d_{\lambda_Q}(v_r v_s) + \lambda_Q(v_r v_s)) = (l_1 + \widetilde{\mu}_Q(v_r v_s), l_2 + \lambda_Q(v_r v_s))$  Since,  $\widetilde{\mu}_Q(v_i v_j) \neq \widetilde{\mu}_Q(v_r v_s)$  and  $\lambda_Q(v_i v_j) \neq \lambda_Q(v_r v_s)$ , we have  $td_G(v_i v_j) \neq td_G(v_r v_s)$ . Hence,  $G$  is not a totally edge regular that is contradiction to our assumption. Therefore,  $(\widetilde{\mu}_P, \lambda_Q)$  is a constant function. Similarly we can show that  $(\widetilde{\mu}_P, \lambda_Q)$  is a constant function, when  $G$  is a totally edge regular cubic graph. □

**Theorem 3.5.** Let  $G = (V, E)$  be a cubic graph on a  $k$ -regular crisp graph  $G^*$ . Then,  $(t_B, f_B)$  is a constant if and only if  $G$  is both regular and edge regular cubic graph.

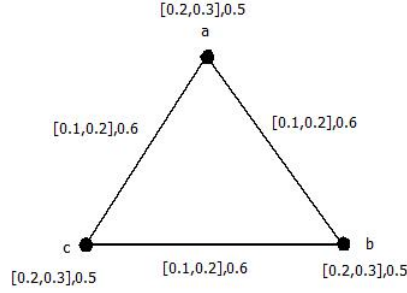


FIGURE 3. Strongly regular cubic graph

*Proof.* Let  $G = (V, E)$  be a cubic graph on  $G^*$  and let  $G^*$  be a  $k$ -regular crisp graph. Assume that  $\widetilde{\mu}_Q$  and  $\lambda_Q$  are constant functions, i.e.,  $\widetilde{\mu}_Q(v_i v_j) = c$  and  $\lambda_Q(v_i v_j) = t$ , for all  $v_i v_j \in E$ , where  $c, t$  are constants. By definition of degree of a vertex we have  $d_G(v_i) = (d_{\widetilde{\mu}_P}(v_i), d_{\lambda_Q}(v_i)) = (\sum_{v_i v_j \in E} \widetilde{\mu}_Q(v_i v_j), \sum_{v_i v_j \in E} \lambda_Q(v_i v_j)) = (\sum_{v_i v_j \in E} c, \sum_{v_i v_j \in E} t)$  for all  $v_i \in V$ . Hence,  $d_G(v_i) = (kc, kt)$ . Therefore,  $G$  is regular cubic graph. Now,  $td_G(v_i v_j) = (td_{\widetilde{\mu}_Q}(v_i v_j), td_{\lambda_Q}(v_i v_j))$ , where  $td_i(v_i v_j) = \sum_{k \neq j} \widetilde{\mu}_Q(v_i v_k) + \sum_{k \neq i} \widetilde{\mu}_P(v_k v_j) + \widetilde{\mu}_Q(v_i v_j) = \sum_{k \neq j} c + \sum_{k \neq i} c + c = c(k-1) + c(k-1) + c = c(2k-1)$ . Similarly,  $td_{\lambda_Q}(v_i v_j) = t(2k-1)$ , for all  $v_i v_j \in E$ . Hence,  $G$  is also totally edge regular cubic graph.

Conversely, assume that  $G$  is both regular and edge regular cubic graph. We prove that  $(\widetilde{\mu}_Q, \lambda_Q)$  is a constant function. Since  $G$  is regular,  $d_G(v_i) = (c_1, c_2)$ , for all  $v_i \in V$ . Also,  $G$  is totally edge regular. Hence  $td_G(v_i v_j) = (t_1, t_2)$ , for all  $v_i v_j \in E$ . By definition of totally edge degree we have  $td_G(v_i v_j) = (td_{\widetilde{\mu}_Q}(v_i v_j), td_{\lambda_Q}(v_i v_j))$ , where  $td_G(v_i v_j) = d_{\widetilde{\mu}_P}(v_i) + d_{\widetilde{\mu}_P}(v_j) - \widetilde{\mu}_Q(v_i v_j)$ , for all  $v_i v_j \in E$ ,  $t_1 = c_1 + c_2 - \widetilde{\mu}_Q(v_i v_j)$ . So,  $\widetilde{\mu}_Q(v_i v_j) = 2c_1 - t_1$ . Similarly we have  $\lambda_Q(v_i v_j) = 2c_2 - t_2$ , for all  $v_i v_j \in E$ . Hence,  $(\widetilde{\mu}_P, \lambda_Q)$  is a constant function.  $\square$

**Definition 3.3.** A cubic graph  $G = (V = \{v_1, v_2, \dots, v_n\}, E)$ , is said to be strongly regular, if it satisfies the following axioms:

- (i)  $G$  is  $k = (k_1, k_2)$ -regular cubic graph
- (ii) The sum of membership values and non-membership values of the common neighborhood vertices of any pair of adjacent vertices and non-adjacent vertices  $v_i, v_j$  of  $G$  has the same weight and is denoted by  $\lambda = (\widetilde{\lambda}_1, \lambda_2)$ ,  $\delta = (\widetilde{\delta}_1, \delta_2)$ , respectively.

**Note 3.1.** Any strongly cubic graph  $G$  is denoted by  $G = (n, k, \lambda, \delta)$ .

**Example 3.3.** Consider the cubic graph  $G = (V, E)$  in Figure 3, where  $V = \{a, b, c\}$  and  $\widetilde{\mu}_Q(a, b) = \widetilde{\mu}_Q(b, a) = \widetilde{\mu}_Q(a, c) = \widetilde{\mu}_Q(c, a) = \widetilde{\mu}_Q(b, c) = \widetilde{\mu}_Q(c, b) = (k_1, k_2) = ([0.2, 0.4], 1.2)$ ,  $\lambda = (\widetilde{\lambda}_1, \lambda_2, \lambda^*) = [(0.4, 0.6), 1.2)$ ,  $\delta = (\widetilde{\delta}_1, \delta_2) = [(0, 0), 0]$ . Hence,  $G$  is a strongly regular cubic graph.

**Theorem 3.6.** If  $G = (V, E)$  is a complete cubic graph with  $(t_A, f_A)$  and  $(t_B, f_B)$  as constant functions, then  $G$  is a strongly regular cubic graph.



*Proof.* Let  $G = (V, E)$  be a complete cubic graph where  $V = \{v_1, v_2, \dots, v_n\}$ . Since  $\widetilde{\mu}_P(v_i)$ ,  $\lambda_Q(v_i)$ ,  $\widetilde{\mu}_P(v_i v_j)$  and  $\lambda_Q(v_i v_j)$  are constant functions, hence,  $\widetilde{\mu}_P(v_i) = r$ ,  $\lambda_Q(v_i) = s$ , for all  $v_i \in V$  and  $\widetilde{\mu}_P(v_i v_j) = c$  and  $\lambda_Q(v_i v_j) = t$ , for all  $v_i v_j \in E$  where  $r, s, c, t$  are constants. To prove that  $G$  is a strongly regular cubic graph, we have to show that  $G$  is  $k = (k_1, k_2)$ -regular cubic graph and the adjacent vertices have the same common neighborhood  $\lambda = (\lambda_1, \lambda_2)$  and non-adjacent vertices have the same common neighborhood  $\delta = (\delta_1, \delta_2)$ . Now, Since  $G$  is complete;  $d_G(v_i) = (d_{\widetilde{\mu}_P}(v_i), d_{\lambda_Q}(v_i)) = (\sum_{v_i v_j \in E} \widetilde{\mu}_P(v_i v_j), \sum_{v_i v_j \in E} \lambda_Q(v_i v_j)) = ((n-1)c, (n-1)t)$  Hence,  $G$  is an  $((n-1)c, (n-1)t)$ -regular cubic graph. Now, since  $G$  is complete cubic graph, the sum of membership values and non-membership values of common neighborhood vertices of any pair of adjacent vertices  $\lambda = ((n-2)r, (n-2)s)$  are the same and the sum of membership values and non-membership values of common neighborhood vertices of any pair of non-adjacent vertices  $\delta = 0$  are the same.  $\square$

**Remark 3.2.** If  $G$  is a strongly regular disconnected cubic graph then,  $\delta = 0$ .

**Definition 3.4.** A cubic graph  $G = (V, E)$  is said to be a biregular cubic graph if it satisfies the following axioms:

- (i)  $G$  is  $k = (k_1, k_2)$ -regular cubic graph.
- (ii)  $V = V_1 \cup V_2$  be the bipartition of  $V$  and every vertex in  $V_1$  has the same neighborhood degree  $M = (M_1, M_1)$  and every vertex in  $V_2$  has the same neighborhood degree  $N = (N_1, N_2)$ , where  $M$  and  $N$  are constants.

**Example 3.4.** Consider a cubic graph  $G = (V, E)$  in Figure 4, where  $V = \{u_1([0.25, 0.35], 0.5), u_2([0.35, 0.45], 0.6), u_3([0.25, 0.35], 0.5), u_4([0.35, 0.45], 0.6), u_5([0.35, 0.45], 0.6), u_6([0.25, 0.35], 0.5), u_7([0.35, 0.45], 0.6), u_8([0.25, 0.35], 0.5)\}$  and  $E = \{u_1 u_2, u_1 u_4, u_1 u_5, u_2 u_6, u_2 u_3, u_3 u_4, u_3 u_7, u_4 u_8, u_5 u_6, u_5 u_8, u_6 u_7, u_7 u_8\}$ . The membership values of the edges  $(u_1, u_2)$ ,  $(u_3, u_4)$ ,  $(u_5, u_8)$ ,  $(u_6, u_7)$  is  $([0.05, 0.15], 0.6)$  and that of  $(u_1, u_4)$ ,  $(u_2, u_3)$ ,  $(u_5, u_6)$ ,  $(u_7, u_8)$  is  $([0.15, 0.25], 0.7)$  and that of  $(u_1, u_5)$ ,  $(u_2, u_6)$ ,  $(u_4, u_8)$ ,  $(u_3, u_7)$  is  $([0.25, 0.35], 0.6)$ . Then  $k = (k_1, k_2) = ([0.55, 0.65], 1.9)$ ,  $V_1 = \{u_1, u_3, u_6, u_8\}$ ,  $V_2 = \{u_2, u_4, u_5, u_7\}$ ,  $M = (M_1, M_2) = ([1.15, 1.25], 1.8)$  and  $N = (N_1, N_2) = ([0.85, 0.95], 1.5)$ .

**Theorem 3.7.** If  $G = (V, E)$  is a strongly regular cubic graph which is strong then,  $\overline{G}$  is a  $(k_1, k_2)$ -regular.

*Proof.* Let  $G = (V, E)$  be a strongly regular cubic graph. Then by definition,  $G$  is a  $(k_1, k_2)$ -regular. Since  $G$  is strong, we have

$$\overline{\mu}_Q(v_i v_j) = \begin{cases} 0 & \text{for all } v_i v_j \in E \\ \min(\widetilde{\mu}_P(v_i), \widetilde{\mu}_P(v_j)) & \text{for all } v_i v_j \notin E \end{cases}$$

$$\overline{\lambda}_Q(v_i v_j) = \begin{cases} 0 & \text{for all } v_i v_j \in E \\ \max(\lambda_P(v_i), \lambda_P(v_j)) & \text{for all } v_i v_j \notin E. \end{cases}$$

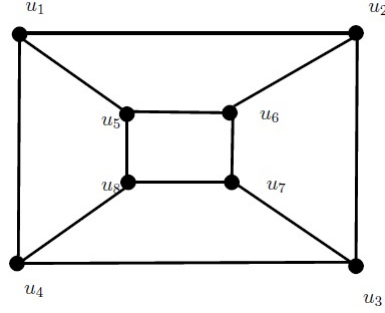


FIGURE 4. Biregular cubic graph

Now, since  $G$  is strong, the degree of a vertex  $v_i$  in  $\bar{G}$  is  $d_{\bar{G}}(v_i) = (d_{\bar{\mu}_P}(v_i), d_{\bar{\lambda}_P}(v_i))$ , where  $d_{\bar{\mu}_P}(v_i) = \sum_{v_i \neq v_j} \bar{\mu}_Q(v_i v_j) = \sum_{v_i \neq v_j} \bar{\mu}_P(v_i) \wedge \bar{\mu}_P(v_j) = k_1$ ,  $\forall v_i \in V$ . Hence,  $d_{\bar{G}}(v_i) = (k_1, k_2)$ , for all  $v_i \in V$ . So,  $\bar{G}$  is a  $(k_1, k_2)$ -regular cubic graph.  $\square$

**Theorem 3.8.** *Let  $G = (V, E)$  be a strong cubic graph. Then  $G$  is a strongly regular if and only if  $\bar{G}$  is a strongly regular.*

*Proof.* Assume that  $G = (V, E)$  is a strongly regular cubic graph. Then  $G$  is  $(k_1, k_2)$ -regular and the adjacent vertices and the non-adjacent vertices have the same common neighborhood  $\lambda = (\lambda_1, \lambda_2)$  and  $\delta = (\delta_1, \delta_2)$ , respectively. We have to prove that  $\bar{G}$  is a strongly regular cubic graph. If  $G$  is strongly regular cubic graph which is strong then  $\bar{G}$  is a  $(k_1, k_2)$ -regular cubic graph by Theorem 3.7. Next, let  $S_1$  and  $S_2$  be the sets of all adjacent vertices and non-adjacent vertices of  $G$ . That is,  $S_1 = \{v_i v_j \mid v_i v_j \in E\}$ , where  $v_i$  and  $v_j$  have same common neighborhood  $\lambda = (\lambda_1, \lambda_2)$  and  $S_2 = \{v_i v_j \mid v_i v_j \notin E\}$ , where  $v_i$  and  $v_j$  have same common neighborhood  $\delta = (\delta_1, \delta_2)$ . Then,  $\bar{S}_1 = \{v_i v_j \mid v_i v_j \in \bar{E}\}$ , where  $v_i$  and  $v_j$  have same common neighborhood  $\delta = (\delta_1, \delta_2)$  and  $\bar{S}_2 = \{v_i v_j \mid v_i v_j \notin \bar{E}\}$ , where  $v_i$  and  $v_j$  have same common neighborhood  $\lambda = (\lambda_1, \lambda_2)$ . Which implies  $\bar{G}$  is a strongly regular. Similarly we can prove the converse.  $\square$

**Theorem 3.9.** *A strongly regular cubic graph  $G$  is a biregular cubic graph if the adjacent vertices have the same common neighborhood  $\lambda = (\lambda_1, \lambda_2) \neq 0$  and the non-adjacent vertices have the same common neighborhood  $\delta = (\delta_1, \delta_2) \neq 0$ .*

*Proof.* Let  $G = (V, E)$  be a strongly regular cubic graph. Then we have  $d(v_i) = (k_1, k_2)$ , for all  $v_i \in V$ . Assume that the adjacent vertices have the same common neighborhood  $\delta = (\delta_1, \delta_2) \neq 0$ . Let  $S$  be the sets of all non-adjacent vertices. That is  $S = \{v_i v_j \mid v_i$  is not adjacent to  $v_j, i \neq j, v_i, v_j \in V\}$ . Now the vertex partition of  $G$  is  $V_1 = \{v_i \mid v_i \in S\}$  and  $V_2 = \{v_j \mid v_j \in S\}$ . Then  $V_1$  and  $V_2$  have the same neighborhood degree, since  $G$  is a strongly regular. Hence,  $G$  is a bi-regular cubic graph.  $\square$

**Definition 3.5.** (i) *If the underlying graph  $G^*$  is an edge regular graph, then  $G$  is said to be a partially edge regular cubic graph.*

(ii) If  $G$  is both edge regular and partially edge regular cubic graph, then  $G$  is said to be a full edge regular cubic graph.

**Theorem 3.10.** Let  $G$  be a cubic graph on  $G^*$  such that  $(t_B, f_B)$  is a constant function. If  $G$  is full regular, then  $G$  is full edge regular cubic graph.

*Proof.* Let  $(\widetilde{\mu}_Q, \lambda_Q)$  be a constant function. Then,  $\widetilde{\mu}_Q(v_i v_j) = c_1$  and  $\lambda_Q(v_i v_j) = c_2$ , for every  $v_i v_j \in E$  where  $c_1$  and  $c_2$  are constants. Suppose that  $G$  is full regular cubic graph then  $d_G(v_i) = (k_1, k_2) = k$  and  $d_{G^*}(v_i) = r$ , for all  $v_i \in V$ , where  $k$  and  $r$  are constants.  $d_{G^*}(v_i v_j) = d_{G^*}(v_i) + d_{G^*}(v_j) - 2 = 2r - 2 = \text{constant}$ . Hence,  $G^*$  is an edge regular graph. Now, since  $G$  is regular,  $d_G(v_i v_j) = (d_{\widetilde{\mu}_Q}(v_i v_j), d_{\lambda_Q}(v_i v_j))$ , for all  $v_i v_j \in E$  where  $d_{\widetilde{\mu}_Q}(v_i v_j) = d_{\widetilde{\mu}_P}(v_i) + d_{\widetilde{\mu}_P}(v_j) - 2\widetilde{\mu}_Q(v_i v_j) = k_1 + k_1 - 2c_1 = 2k_1 - 2c_1 = \text{constant}$ . Similarly, for all  $v_i v_j \in E$ ,  $d_{\lambda_Q}(v_i v_j) = 2k_2 - 2c_2 = \text{constant}$ . Hence,  $G$  is an edge regular cubic graph. Therefore,  $G$  is full edge regular cubic graph.  $\square$

**Theorem 3.11.** Let  $G$  be a  $t$ -totally edge regular cubic graph and  $t_1$ -partially edge regular cubic graph. Then  $S(G) = \frac{qt}{1+t_1}$ , where  $q = |E|$ .

*Proof.* The size of cubic graph  $G$  is

$$S(G) = \left( \sum_{v_i v_j \in E} \widetilde{\mu}_Q(v_i v_j), \sum_{v_i v_j \in E} \lambda_Q(v_i v_j) \right).$$

Since  $G$  is  $t$ -totally edge regular cubic graph i.e.,  $td_G(v_i v_j) = t$  and  $G^*$  is  $t_1$ -partially edge regular cubic graph i.e.  $d_{G^*}(v_i v_j) = t_1$ . Thus,  $\sum td_G(v_i v_j) = \left( \sum_{v_i v_j \in E} d_{G^*}(v_i v_j) \widetilde{\mu}_Q(v_i v_j), \sum_{v_i v_j \in E} d_{G^*}(v_i v_j) \lambda_Q(v_i v_j) \right) + S(G)$ .  $qt = t_1 S(G) + S(G)$ . Hence,  $S(G) = \frac{qt}{1+t_1}$ .  $\square$

#### 4. CUBIC DIGRAPH IN TRAVEL TIME

In modern age, planning efficient routes is essential for industry and business, with applications as varied as product distribution, laying new fiber optic lines for broadband internet, and suggesting new friends within social network websites such as Facebook. When we visit a website like Google Maps and looking for directions from one city to another city in USA, we are usually asking for a shortest path between the two cities. These computer applications use representations of the road maps as graphs, with estimated travel times as edge weights. The travel time is a function of traffic density on the road or the length of the road. The traffic density is a fuzzy, while the length of a road is a crisp quantity. In a road network, crossings are represented by vertices, roads by edges and traffic density on the road is usually calculated between adjacent crossings. These factors can be represented as a cubic set. Any model of a road network can be represented as a cubic digraph  $D = (C, R)$ , where  $C$  is a cubic set of crossings(vertices) at which the traffic density is calculated and connectivity conditions as truth-membership degree with intervals  $\widetilde{\mu}_P(x)$  and falsity membership degree  $\lambda_P(x)$

$$C = \left\{ (a, [0.2, 0.3], 0.4), (b, [0.3, 0.5], 0.6), (c, [0.5, 0.7], 0.8), (d, [0.4, 0.6], 0.7), (e, [0.2, 0.5], 0.6), (f, [0.3, 0.4], 0.5), (g, [0.4, 0.6], 0.7) \right\}$$

and  $R$  is a cubic set of roads (edges) between crossings whose truth membership degree

$\widetilde{\mu}_P(x)$  and false membership degree  $\lambda_P(x)$  can be calculated as:  
 $\widetilde{\mu}_Q(xy) \leq \min\{\widetilde{\mu}_P(x), \widetilde{\mu}_P(y)\}$   
 $\lambda_P(xy) \geq \max\{\lambda_P(x), \lambda_P(y)\}$  for all  $xy \in E$ .

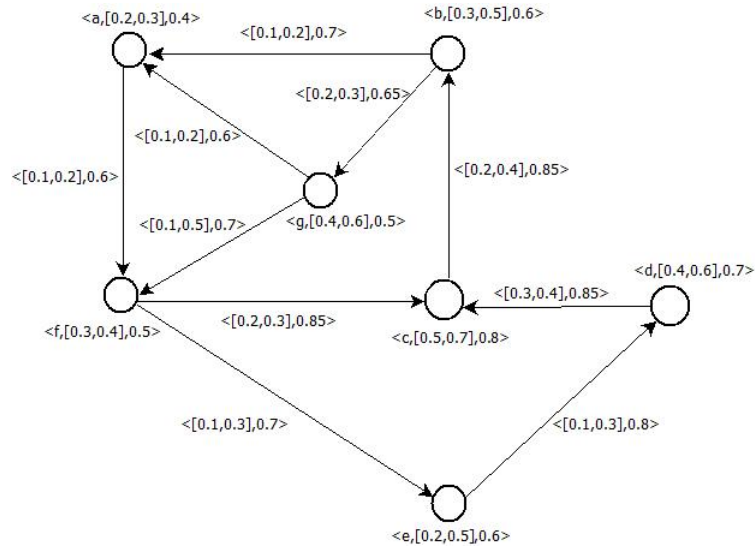


FIGURE 5. Cubic digraph of a road network

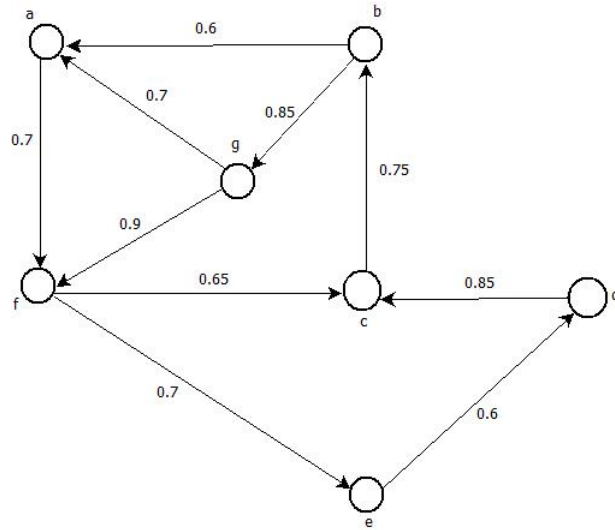


FIGURE 6. Weighted digraph of a road network

The cubic digraph  $D = (C, R)$  of the travel time is given in Figure 5. The cubic out neighborhood are given in Table 4.

The final weights on edges can be calculated by finding the score function of cubic edges as

$s_i = (\widetilde{\mu}_R)_i + 1 - (\lambda_R)_i$ . The final weighted digraph given in Figure 6 which can be used for finding the shortest or optimal path between two locations (vertices) by any of the known methods, like Floyd’s algorithm, Dijkstra’s and A star algorithm. Weighted relations are given in Table 4.

Crossings	$N^+$ (crossings)
a	{f,([0.1,0.2],0.6)}
b	{a,([0.1,0.2],0.7), g,([0.2,0.3],0.65)}
c	{b,([0.2,0.4],0.85)}
d	{c,([0.3,0.4],0.85)}
e	{d,([0.1,0.3],0.8)}
f	{c,([0.2,0.3],0.85), e,([0.1,0.3],0.7)}
g	{a,([0.1,0.2],0.6), f,([0.1,0.5],0.7)}

Crossings	$N^+$ (crossings)
a	{f,0.7}
b	{a,0.6 , g,0.85}
c	{b,0.75}
d	{c,0.85}
e	{d,0.6}
f	{c,0.65, e,0.7}
g	{a,0.7, f,0.9}

The following algorithm generates the weighted digraph, WR, for given cubic digraph and resolves to obtain the optimal path from a source vertex to destination vertex.

### 5. CONCLUSION

Fuzzy graph theory has numerous applications in modern science and technology, especially in the fields of operations research, neural networks, and decision making. Since the cubic models give more precision, flexibility and compatibility to the system as compared to the classical and fuzzy models, in this paper the definition of partial edge regular and fully edge regular cubic graph are given and some properties of edge regular cubic graph are studied. Also, we have introduced the condition under which edge regular cubic graph and totally edge regular cubic graph are equivalent. In our future work, we are going to extend the properties of strongly edge regular cubic graph in matrix representation.

**Algorithm 1** Cubic digraphs in road networks

---

```

1: void cubic shortest path()
2: C = cubic set of crossings;
3: number of crossings = count(C);
4: R = Empty cubic set of roads;
5: for(int c = 0; c < number of crossings ; c++){
6:   for(int c' = 0; c' < number of crossings ; c'++){
7:     if (C(x) is adjacent to C(y)){
8:        $\overline{\mu}_{PR}cc' \leq \min(\overline{\mu}_Q(c), \overline{\mu}_Q(c'))$ ;
9:        $\lambda_{PCC'} \geq \max(\lambda_Q(c), \lambda_Q(c'))$ ;
10:    }
11:  }
12: }
13: R = cubic set of edges;
14: R = cubic relation (Adjacency matrix of  $F \times F$ );
15: WG = Weighted relation (Adjacency matrix of  $F \times F$ );
16: no. of Edges = Count(R);
17: for(int i=0; i < no. of Edges ; i++){
18:    $s_i = (\overline{\mu}_R)_i + 1 - (\lambda_R)_i$ ;
19:   c = Adjacent from Node of  $R_i$ ;
20:   c' = Adjacent from Node of  $R_i$ ;
21:    $WR_{cc'} = s_i$ ;
22: }
23: print WR
24: Calculated optimal path using WR
25: }

```

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